

SRT_2^2 does not imply COH under ω -models

Ludovic Patey
Joint work with Benoit Monin



REVERSE MATHEMATICS

Foundational program that seeks to determine the **optimal** axioms of **ordinary** mathematics.

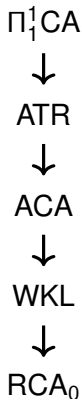
$$\text{RCA}_0 \vdash A \leftrightarrow T$$

in a very weak theory RCA_0
capturing **computable mathematics**

REVERSE MATHEMATICS

Mathematics are
 computationally
 very structured

Almost every theorem is
 empirically equivalent to one
 among five big subsystems.

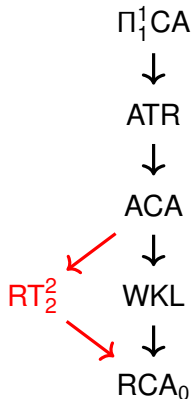


REVERSE MATHEMATICS

Mathematics are
computationally
very structured

Almost every theorem is empirically equivalent to one among five big subsystems.

Except for Ramsey's theory...



RAMSEY'S THEOREM

$[X]^n$ is the set of **unordered n -tuples** of elements of X

A **k -coloring** of $[X]^n$ is a map $f : [X]^n \rightarrow k$

A set $H \subseteq X$ is **homogeneous** for f if $|f([H]^n)| = 1$.

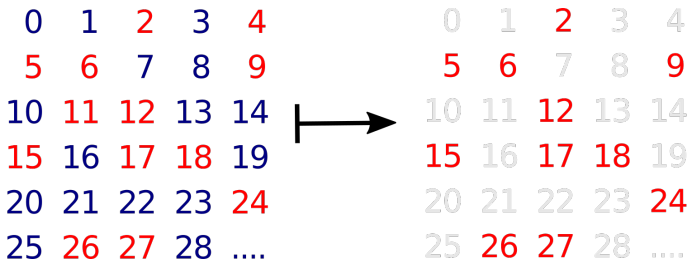
RT _{k} ^{n}

Every k -coloring of $[\mathbb{N}]^n$ admits
an infinite homogeneous set.

PIGEONHOLE PRINCIPLE

RT_k^1

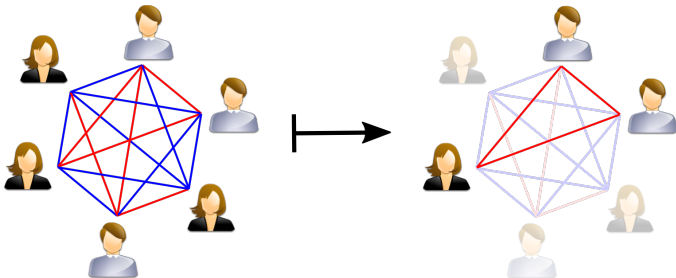
Every k -partition of \mathbb{N} admits an infinite part.



RAMSEY'S THEOREM FOR PAIRS

RT_k^2

Every k -coloring of the infinite clique admits an infinite monochromatic subclique.



An infinite set C is \vec{R} -cohesive for some sets R_0, R_1, \dots if for every i , either $C \subseteq^* R_i$ or $C \subseteq^* \overline{R}_i$.

COH : Every collection of sets has a cohesive set.

A coloring $f : [\omega]^2 \rightarrow 2$ is **stable** if $\lim_y f(x, y)$ exists for every x .

SRT₂² : Every stable coloring of pairs admits an infinite homogeneous set.

$$\text{RCA}_0 \vdash \text{RT}_2^2 \leftrightarrow \text{COH} \wedge \text{SRT}_2^2$$

(Cholak, Jockusch and Slaman)

- ▶ Given $f : [\mathbb{N}]^2 \rightarrow 2$, define $\langle R_x : x \in \mathbb{N} \rangle$ by $R_x = \{y : f(x, y) = 1\}$
- ▶ By COH, there is an \vec{R} -cohesive set $C = \{x_0 < x_1 < \dots\}$
- ▶ $f : [C]^2 \rightarrow 2$ is stable

$$\text{RCA}_0 \vdash \text{RT}_2^2 \leftrightarrow \text{COH} \wedge \text{SRT}_2^2$$

(Cholak, Jockusch and Slaman)

Thm (Hirschfeldt, Jockusch, Kjos-Hanssen, Lempp, and Slaman)

$$\text{RCA}_0 \not\vdash \text{COH} \rightarrow \text{SRT}_2^2$$

Thm (Chong, Slaman and Yang)

$$\text{RCA}_0 \not\vdash \text{SRT}_2^2 \rightarrow \text{COH}$$

Using a [non-standard model](#) containing only low sets.

Thm (Monin and Patey)

$RCA_0 \not\leq SRT_2^2 \rightarrow COH$ on ω -models

Thm (Monin and Patey)

$RCA_0 \not\vdash SRT_2^2 \rightarrow COH$ on ω -models

Jump cone avoidance

RT^2 admits low_2 solutions
CJS trick: largeness



Jump PA avoidance

SRT^2 does not imply COH
Monin's trick: largeness on product spaces



Cone avoidance

Seetapun's theorem
CJS trick: overapproximations



PA avoidance

Liu's theorem
Liu's trick: valuations

Reformulate the problem using
computability-theoretic
versions of the statements

COH admits a **universal** instance:
the primitive recursive sets

A set is **p-cohesive** if it is cohesive for the p.r. sets

Thm (Jockusch and Stephan)

A set is p-cohesive iff its jump is PA over \emptyset'

Thm (Jockusch and Stephan)

For every computable sequence of sets \vec{R} and every p-cohesive set C , C computes an \vec{R} -cohesive set.

SRT₂² can be seen as a Δ_2^0 instance of
the pigeonhole principle

- ▶ Given a stable computable coloring $f : [\omega]^2 \rightarrow 2$
- ▶ Let $A = \{x : \lim_y f(x, y) = 1\}$
- ▶ Every infinite set $H \subseteq A$ or $H \subseteq \bar{A}$ computes an infinite f -homogeneous set.

Thm (Monin and Patey)

For every Δ_2^0 set A , there is an infinite set $H \subseteq A$ or $H \subseteq \bar{A}$ whose jump is not of PA degree over \emptyset' .

The proof relativizes as follows

Thm (Monin and Patey)

Fix a set Z whose jump is not of PA degree over \emptyset' . For every $\Delta_2^0(Z)$ set A , there is an infinite set $H \subseteq A$ or $H \subseteq \bar{A}$ such that $(H \oplus Z)'$ is not of PA degree over \emptyset' .

LAYER 1: CONE AVOIDANCE

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STRONG CONE AVOIDANCE

A problem P admits **strong cone avoidance** if for every non-computable set C and every P -instance X , there is a solution Y to X such that $C \not\leq_T Y$.

Thm (Dzhafarov, Jockusch)

RT_2^1 admits strong cone avoidance

Forcing question

$$c \text{ ?} \Vdash \varphi(G)$$

where $c \in \mathbb{P}$ and $\varphi(G)$ is Σ_1^0

Lem

Let $c \in \mathbb{P}$ and $\varphi(G)$ be a Σ_1^0 formula.

- (a) If $c \text{ ?} \Vdash \varphi(G)$, then $d \Vdash \varphi(G)$ for some $d \leq c$;
- (b) If $c \text{ ?} \not\Vdash \varphi(G)$, then $d \Vdash \neg\varphi(G)$ for some $d \leq c$.

Suppose $c \not\vdash \varphi(G)$ is uniformly Σ_1^0 whenever $\varphi(G)$ is Σ_1^0

Lem

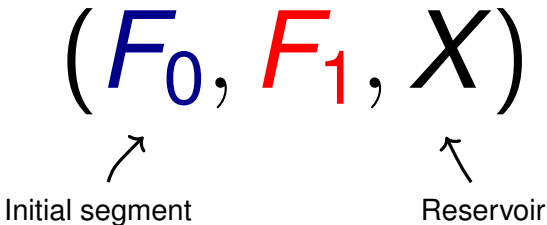
For every non-computable set C and Turing functional Φ_e , the following set is dense in (\mathbb{P}, \leq) .

$$D = \{c \in \mathbb{P} : c \Vdash \Phi_e^G \neq C\}$$

Given $c \in \mathbb{P}$, define the Σ_1^0 set $W = \{(x, v) : c \not\vdash \Phi_e^G(x) \downarrow = v\}$

There is some x such that $(x, 1 - C(x)) \in W$ or $(x, C(x)) \notin W$ otherwise we compute C .

PIGEONHOLE FORCING



- ▶ F_i is **finite**, X is **infinite**, $\max F_i < \min X$ (Mathias condition)
- ▶ $X \in \mathcal{M} \models \text{WKL}$ (Weakness property)
- ▶ $F_i \subseteq A_i$ (Combinatorics)

FORCING RELATION

$$(F_0, F_1, X) \Vdash_i \varphi(G_i)$$

$$\Sigma_1^0 (F_0, F_1, X) \Vdash_i (\exists x)\varphi(G_i, x) \quad \text{if } (\exists w \in \omega)\varphi(F_i, w)$$

$$\Pi_1^0 (F_0, F_1, X) \Vdash_i (\forall x)\varphi(G_i, x) \quad \text{if } (\forall E \subseteq X)(\forall w)\varphi(F_i \cup E, w)$$

FORCING QUESTION

$$(F_0, F_1, X) \text{ ?} \Vdash \varphi_0(G_0) \vee \varphi_1(G_1)$$

Lem

Let $c \in \mathbb{P}$ and $\varphi_0(G), \varphi_1(G)$ be a Σ_1^0 formulas.

- (a) If $c \text{ ?} \Vdash \varphi_0(G_0) \vee \varphi_1(G_1)$, then $d \Vdash_i \varphi_i(G_i)$
- (b) If $c \text{ ?} \not\Vdash \varphi_0(G_0) \vee \varphi_1(G_1)$, then $d \Vdash_i \neg \varphi_i(G_i)$

for some $d \leq c$ and $i < 2$.

CJS TRICK: OVER-APPROXIMATIONS

$$(F_0, F_1, X) ?\vdash \varphi_0(G_0) \vee \varphi_1(G_1)$$

is the $\Sigma_1^{0,X}$ formula

$$(\forall B_0 \sqcup B_1 = X)(\exists i < 2)(\exists E \subseteq B_i)\varphi_i(F_i \cup E)$$

CJS TRICK: OVER-APPROXIMATIONS

- Case 1: $(F_0, F_1, X) ?\vdash \varphi_0(G_0) \vee \varphi_1(G_1)$

Letting $B_i = A_i$, there is an extension $d \leq c$ such that

$$d \Vdash^0 \varphi_1(G_0) \quad \text{or} \quad d \Vdash^1 \varphi_1(G_1)$$

- Case 2: $(F_0, F_1, X) ?\not\vdash \varphi_0(G_0) \vee \varphi_1(G_1)$

The class \mathcal{C} of all $B_0 \sqcup B_1 = \mathbb{N}$ such that

$$(\forall i < 2)(\forall E_i \subseteq X \cap B_i) \Phi_{e_i}^{F_i \cup E_i}(x) \neq n$$

is a non-empty $\Pi_1^{0,X}$ class. Pick $B_0 \sqcup B_1 \in \mathcal{C} \cap \mathcal{M}$.

$$(F_0, F_1, X \cap B_i) \Vdash_i \neg \varphi_i(G_i)$$

LAYER 2: PA AVOIDANCE

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Thm (Liu)

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Forcing question

$$c \text{ ?} \Vdash \varphi(G)$$

where $c \in \mathbb{P}$ and $\varphi(G)$ is Σ_1^0

Lem

Let $c \in \mathbb{P}$ and $\varphi(G)$ and $\psi(G)$ be a Σ_1^0 formulas. If $c \text{ ?} \not\Vdash \varphi(G)$ and $c \text{ ?} \not\Vdash \psi(G)$, then $d \Vdash \neg\varphi(G) \wedge \neg\psi(G)$ for some $d \leq c$.

Suppose $c \not\vdash \varphi(G)$ is uniformly Σ_1^0 whenever $\varphi(G)$ is Σ_1^0

Lem

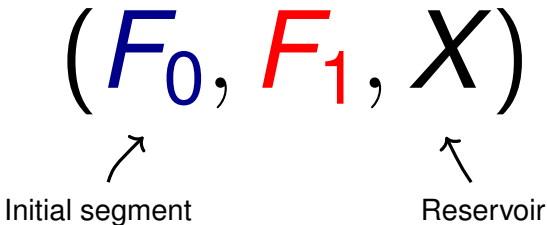
For every Turing functional Φ_e , the following set is dense in (\mathbb{P}, \leq) .

$$D = \{c \in \mathbb{P} : c \Vdash \Phi_e^G \text{ is not a } DNC_2 \text{ function} \}$$

Given $c \in \mathbb{P}$, define the Σ_1^0 set $W = \{(x, v) : c \not\vdash \Phi_e^G(x) \downarrow = v\}$

There is some x such that $(x, \Phi_x(x) \downarrow) \in W$ or $(x, 0), (x, 1) \notin W$ otherwise we compute a DNC_2 function.

PIGEONHOLE FORCING



- ▶ F_i is **finite**, X is **infinite**, $\max F_i < \min X$ (Mathias condition)
- ▶ $X \in \mathcal{M} \models \text{RWKL}$ (Weakness property)
- ▶ $F_i \subseteq A_i$ (Combinatorics)

FORCING QUESTION

$$(F_0, F_1, X) \text{ ?} \vdash \varphi_0(G_0) \vee \varphi_1(G_1)$$

Lem

Let $c \in \mathbb{P}$ and $\varphi_0(G), \varphi_1(G), \psi_0(G), \psi_1(G)$ be a Σ_1^0 formulas.
If $c \text{ ?} \not\vdash \varphi_0(G_0) \vee \varphi_1(G_1)$ and $c \text{ ?} \not\vdash \psi_0(G_0) \vee \psi_1(G_1)$, then

$$d \Vdash_i \neg \varphi_i(G_i) \text{ and } d \Vdash^j \neg \psi_j(G_j)$$

for some $d \leq c$ and $i, j < 2$.

Bad case: $i \neq j$

LIU'S TRICK: VALUATIONS

A **valuation** is a finite function $p \subseteq \omega \rightarrow 2$

A valuation is **correct** if $p(x) = \Phi_x(x) \downarrow$ for every $x \in \text{dom}(p)$

Two valuations p and q are **incompatible** if $p(x) \neq q(x)$ for some $x \in \text{dom}(p) \cap \text{dom}(q)$

Lem (Liu)

Let W be a c.e. set of valuations. Either W contains a correct valuation, or for every $k \in \omega$, there are k pairwise incompatible valuations outside W .

LIU'S TRICK: VALUATIONS

Consider the $\Sigma_1^{0,X}$ set of valuations

$$W = \{p : c \text{ ?}\vdash p \text{ incompatible with either } \Phi_{e_0}^{G_0} \text{ or } \Phi_{e_1}^{G_1}\}$$

Assuming X is not of PA degree, either there is a correct valuation $p \in W$, or three pairwise incorrect valuations $p_0, p_1, p_2 \notin W$.

In the latter case, there is an extension $d \leq c$ and some $i_0, i_1, i_2 < 2$ such that $d \Vdash^{i_s} p_s$ compatible with $\Phi_{e_{i_s}}^{G_{i_s}}$ for each $s < 3$.

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A problem P admits **strong jump cone avoidance** if for every non- Δ_2^0 set C and every P -instance X , there is a solution Y to X such that C is not $\Delta_2^0(Y)$.

Thm (Monin, Patey)

RT_2^1 admits strong jump cone avoidance

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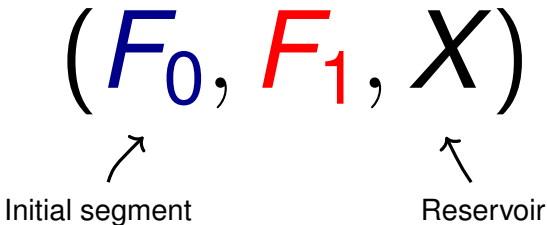
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- ▶ $X \in \mathcal{M} \models \text{WKL}$ (Weakness property)
- ▶ $F_i \subseteq A_i$ (Combinatorics)

FORCING Π_2^0 STATEMENTS

Fix a Δ_0 formula $\psi(G, x, y)$ and a condition c .

$$c \Vdash (\forall x)(\exists y)\psi(G, x, y)$$

$$\equiv$$

$$(\forall x)(\forall d \leq c)(\exists e \leq d)e \Vdash (\exists y)\psi(G, x, y)$$

The quantifier $\forall(E_0, E_1, Y) \leq (F_0, F_1, X)$ is $\Pi_3^0(\mathcal{M})$

CJS TRICK: LARGENESS

- Case 1: $(F_0, F_1, X) ?\vdash \varphi_0(G_0) \vee \varphi_1(G_1)$

Letting $B_i = A_i$, there is an extension $d \leq c$ such that

$$d \Vdash^0 \varphi_1(G_0) \quad \text{or} \quad d \Vdash^1 \varphi_1(G_1)$$

- Case 2: $(F_0, F_1, X) ?\not\vdash \varphi_0(G_0) \vee \varphi_1(G_1)$

The class \mathcal{C} of all $B_0 \sqcup B_1 = \mathbb{N}$ such that

$$(\forall i < 2)(\forall E_i \subseteq X \cap B_i) \Phi_{e_i}^{F_i \cup E_i}(x) \neq n$$

is a non-empty $\Pi_1^{0,X}$ class. Pick $B_0 \sqcup B_1 \in \mathcal{C} \cap \mathcal{M}$.

$$(F_0, F_1, X \cap B_i) \Vdash_i \neg \varphi_i(G_i)$$

CJS TRICK: LARGENESS

The only operations on the reservoirs
are **partitioning** and **trimming**.

Definition

A class $\mathcal{L} \subseteq 2^\omega$ is **large** if it satisfies the following properties:

- (1) For every $X \in \mathcal{L}$ and $Y \supseteq X$, $Y \in \mathcal{L}$
- (2) For every $k \in \omega$ and every $X_0 \cup \dots \cup X_{k-1} = \omega$, there is some $j < k$ such that $X_j \in \mathcal{L}$

Fix a countable coded Scott set $\mathcal{M} = \{X_0, X_1, \dots\}$

Fix a sequence of upward-closed Σ_1^0 classes $\mathcal{U}_0, \mathcal{U}_1, \dots \subseteq 2^\omega$

Given $C \subseteq \omega^2$, let

$$\mathcal{U}_C^{\mathcal{M}} = \bigcap_{(i,e) \in C} \mathcal{U}_e^{X_i}$$

Lem

If $\mathcal{U}_C^{\mathcal{M}}$ is not large, then there is some $X_0 \cup \dots \cup X_{k-1} = \omega$ in \mathcal{M} such that for every $i < k$, $X_i \notin \mathcal{U}_C^{\mathcal{M}}$.

FORCING Π_2^0 FORMULAS

Fix a Δ_1^0 formula $\Phi_e(G, x, y)$

$$(F_0, F_1, X, C) \Vdash_i (\forall x)(\exists y)\neg\Phi_e(G_i, x, y)$$

$$\equiv$$

$$(\forall x)(\forall E \subseteq X \cap A^i)(0, \zeta(e, F_i \cup E, x)) \in C$$

where $\mathcal{U}_{\zeta(e, F, x)}^{X_0} = \{Y : (F, Y) \not\Vdash (\forall y)\Phi_e(G, x, y)\}$

FORCING Π_2^0 FORMULAS

Suppose

$$(F_0, F_1, X, C) \Vdash_i (\forall x)(\exists y)\neg\Phi_e(G_i, x, y)$$

Then

$$(\forall x)(\forall E \subseteq X \cap A^i)(0, \zeta(e, F_i \cup E, x)) \in C$$

Fix some

$$(E_0, E_1, Y, D) \leq (F_0, F_1, X, C)$$

For every x ,

$$\mathcal{U}_D^M \subseteq \mathcal{U}_C^M \subseteq \{Z : (E_i, Z) \not\Vdash (\forall y)\Phi_e(G, x, y)\}$$

So

$$(E_i, Z) \not\Vdash (\forall y)\Phi_e(G, x, y)$$

VALIDITY

A condition (F_0, F_1, X, C) is *i-valid* if $X \cap A^i \in \mathcal{U}_C^M$.

Every condition is either 0-valid or 1-valid

Lem

If (F_0, F_1, X, C) is *i-valid* and $(F_i, X) \not\Vdash (\forall y)\Phi_e(G, y)$ then there is an extension (E_0, E_1, Y, C) with $(E_i, Y) \Vdash (\exists y)\neg\Phi_e(G, y)$

FORCING QUESTION

$$(F_0, F_1, X, C) ?\vdash_i (\exists x)(\forall y)\Phi_e(G_i)$$

$$\equiv$$

$\mathcal{U}_C^M \cap_{x \in \omega, E \subseteq X \cap A^i} \{Z : (F_i \cup E, Z) \not\Vdash (\forall y)\Phi_e(G, x, y)\}$ is not large

Lem

Let $c \in \mathbb{P}$, $i < 2$ and $\varphi(G)$ be a Σ_2^0 formula.

- (a) If $c ?\vdash_i \varphi(G_i)$, then $d \Vdash_i \varphi(G_i)$ for some $d \leq c$;
- (b) If $c ?\not\vdash_i \varphi(G_i)$, then $d \Vdash_i \neg\varphi(G_i)$ for some $d \leq c$.

JUMP CONE AVOIDANCE: SUMMARY

Let \mathcal{F} be a sufficiently generic filter for the pigeonhole forcing.

By the non-disjunctive Σ_2^0 forcing question, **for every $i < 2$** and every $e \in \omega$, there is a condition $c \in \mathcal{F}$ such that $c \Vdash_i \Phi_e^{G'} \neq C$.

For every Σ_2^0 formula $\varphi(G)$, whenever $c \Vdash_i \varphi(G_i)$, then $\varphi(G_i)$ holds.

There is some $i < 2$ such that every condition in \mathcal{F} is i -valid.

Therefore, **there is some $i < 2$** such that for every Π_2^0 formula $\varphi(G)$, whenever $c \Vdash_i \varphi(G_i)$, then $\varphi(G_i)$ actually holds.

LAYER 4: JUMP PA AVOIDANCE

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A problem P admits **jump PA avoidance** if for every computable P -instance X , there is a solution Y to X whose jump is of non-PA degree over \emptyset' .

Thm (Monin and Patey)

SRT_2^2 admits jump PA avoidance

Forcing question

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where $c \in \mathbb{P}$ and $\varphi(G)$ is Σ_2^0

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Let $c \in \mathbb{P}$ and $\varphi(G)$ and $\psi(G)$ be a Σ_2^0 formulas. If $c \text{ ? } \not\Vdash \varphi(G)$ and $c \text{ ? } \not\Vdash \psi(G)$, then $d \Vdash \neg\varphi(G) \wedge \neg\psi(G)$ for some $d \leq c$.

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Lem

For every Turing functional Φ_e , the following set is dense in (\mathbb{P}, \leq) .

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Given $c \in \mathbb{P}$, define the Σ_2^0 set $W = \{(x, v) : c \not\vdash \Phi_e^{G'}(x) \downarrow = v\}$

There is some x such that $(x, \Phi_x(x) \downarrow) \in W$ or $(x, 0), (x, 1) \notin W$ otherwise we \emptyset' -compute a $DNC_2^{\emptyset'}$ function.

(F_0, F_1, X, C)



Initial segment



Reservoir



Largeness

- ▶ F_i is **finite**, $\max F_i < \min X$ (Mathias condition)
- ▶ \mathcal{U}_C^M is large, $\mathcal{U}_C^M \subseteq \mathcal{L}_X$ (Second jump control)
- ▶ $X \in \mathcal{M} \models \text{WKL}$, C is Δ_2^0 (Weakness property)
- ▶ $F_i \subseteq A_i$ (Combinatorics)

FORCING QUESTION

$$(F_0, F_1, X, C) \Vdash_i (\exists x)(\forall y)\Phi_e(G_i)$$

$$\equiv$$

$\mathcal{U}_C^M \cap_{x \in \omega, E \subseteq X \cap A_i} \{Z : (F_i \cup E, Z) \Vdash (\forall y)\Phi_e(G, x, y)\}$ is not large

Suppose $c \Vdash_i \varphi(G_i)$ and $c \Vdash_i \psi(G_i)$

Then there are two sets $D_0, D_1 \subseteq \omega^2$ such that

$$(F_0, F_1, X, D_0) \Vdash \neg \varphi(G_i) \text{ and } (F_0, F_1, X, D_1) \Vdash \neg \psi(G_i)$$

But $\mathcal{U}_{D_0}^M \cap \mathcal{U}_{D_1}^M$ may not be large!

MONIN'S TRICK: LARGENESS IN PRODUCT SPACES

Suppose $c \not\Vdash_i \varphi(G_i)$ and $c \not\Vdash_i \psi(G_i)$

Then there are two sets $D_0, D_1 \subseteq \omega^2$ such that

$$(F_0, F_1, X, D_0) \Vdash \neg\varphi(G_i) \text{ and } (F_0, F_1, X, D_1) \Vdash \neg\psi(G_i)$$

Define a new kind of condition

$$(F_0, F_1, Y_0, Y_1, E)$$

where $Y_0 \subseteq X$, $Y_1 \subseteq X$ and $\mathcal{U}_E^M \subseteq \mathcal{U}_{D_0}^M \times \mathcal{U}_{D_1}^M$

Think of it as $(F_0, F_1, Y_0 \cup Y_1)$ where Y_0 and Y_1 must be dependently large

LARGENESS IN PRODUCT SPACES

Definition

A class $\mathcal{L} \subseteq 2^\omega \times \cdots \times 2^\omega$ is **large** if it satisfies the following properties:

- (1) For every $\langle X_i : i < k \rangle \in \mathcal{L}$ and $Y_i \supseteq X_i$, $\langle Y_i : i < k \rangle \in \mathcal{L}$
- (2) For every $\ell \in \omega$ and every $X_0 \cup \cdots \cup X_{\ell-1} = \omega$, there is some $j_0, \dots, j_{k-1} < \ell$ such that $\langle X_{j_i} : i < k \rangle \in \mathcal{L}$

$$(F_0, F_1, X_0, \dots, X_{k-1}, C)$$


Initial segment



Reservoirs



Largeness

- ▶ F_i is **finite**, $\max F_i < \min(X_0, \dots, X_{k-1})$ (Mathias condition)
- ▶ \mathcal{U}_C^M is large, $\mathcal{U}_C^M \subseteq \mathcal{L}_{X_0} \times \dots \times \mathcal{L}_{X_{k-1}}$ (Second jump control)
- ▶ $X_0, \dots, X_{k-1} \in \mathcal{M} \models \text{WKL}$, C is Δ_2^0 (Weakness property)
- ▶ $F_i \subseteq A_i$ (Combinatorics)

PARTIAL ORDER

$$(E_0, E_1, Y_0, \dots, Y_{nk-1}, D) \leq (F_0, F_1, X_0, \dots, X_{k-1}, C)$$

- ▶ $Y_{ks+j} \subseteq X_j$ for every $s < n$ and $j < k$
- ▶ $F_i \subseteq E_i$ and $E_i \setminus F_i \subseteq X_0 \cup \dots \cup X_{k-1}$
- ▶ for every $\langle Z_{ks+j} : s < n, j < k \rangle \in \mathcal{U}_D^M$, and $s < n$,
 $\langle Z_{ks+j} : j < k \rangle \in \mathcal{U}_C^M$

VALIDITY

A condition $(F_0, F_1, X_0, \dots, X_{k-1}, C)$ is *i*-valid if,

$$\langle X_j \cap A^i : j < k \rangle \in \mathcal{U}_C^M$$

Some conditions are neither 0-valid nor 1-valid!

By largeness, there are some $i_0, \dots, i_{k-1} < 2$ such that

$$\langle X_j \cap A^{i_j} : j < k \rangle \in \mathcal{U}_C^M$$

LIU'S TRICK: VALUATIONS

A **valuation** is a finite function $p \subseteq \omega \rightarrow 2$

A valuation is **\emptyset' -correct** if $p(x) = \Phi_x^{\emptyset'}(x) \downarrow$ for every $x \in \text{dom}(p)$

Two valuations p and q are **incompatible** if $p(x) \neq q(x)$ for some $x \in \text{dom}(p) \cap \text{dom}(q)$

Lem (Liu)

Let W be a Σ_2^0 set of valuations. Either W contains a \emptyset' -correct valuation, or for every $k \in \omega$, there are k pairwise incompatible valuations outside W .

LIU'S TRICK: VALUATIONS

Fix a condition $c = (F_0, F_1, X, C)$

Consider the $\Sigma_1^{0,X}$ set of valuations

$$W = \{p : c \Vdash_i p \text{ incompatible with } \Phi_e^{G'_i}\}$$

Assuming X is low, either there is a \emptyset' -correct valuation $p \in W$, or three pairwise incorrect valuations $p_0, p_1, p_2 \notin W$.

In the latter case, there are three sets $D_0, D_1, D_2 \subseteq \omega^2$ such that $\mathcal{U}_{D_0}^M, \mathcal{U}_{D_1}^M$ and $\mathcal{U}_{D_2}^M$ are large, and for each $s < 3$.

$$(F_0, F_1, X, D_s) \Vdash_i p_s \text{ compatible with } \Phi_e^{G'_i}$$

PARALLEL CONDITIONS

Let

$$\mathcal{U}_D^M = \mathcal{U}_{D_0}^M \times \mathcal{U}_{D_1}^M \times \mathcal{U}_{D_2}^M$$

and define

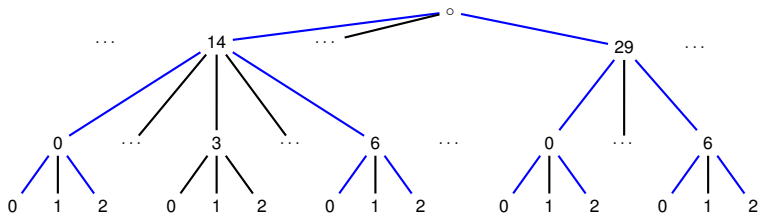
$$d = (F_0^{\{r,s\}}, F_1^{\{r,s\}}, X_0, X_1, X_2, D : r < s < 3)$$

The condition d represents three **parallel** conditions

- ▶ $d^{\{0,1\}} = (F_0^{\{0,1\}}, F_1^{\{0,1\}}, X_0, X_1, \pi_{0,1}(D))$
- ▶ $d^{\{0,2\}} = (F_0^{\{0,2\}}, F_1^{\{0,2\}}, X_0, X_2, \pi_{0,2}(D))$
- ▶ $d^{\{1,2\}} = (F_0^{\{1,2\}}, F_1^{\{1,2\}}, X_1, X_2, \pi_{1,2}(D))$

One among $d^{\{0,1\}}$, $d^{\{0,2\}}$ and $d^{\{1,2\}}$ is i -valid for some $i < 2$

PARALLEL CONDITIONS



FURTHER RESULTS

Thm (Monin and Patey)

For every Δ_2^0 set A , there is an infinite set $H \subseteq A$ or $H \subseteq \bar{A}$ whose jump is not of 2-random degree.

- ▶ Does every set A have an infinite set $H \subseteq A$ or $H \subseteq \bar{A}$ whose jump is not of PA degree relative to \emptyset' ?
- ▶ Does every Δ_2^0 set A have an infinite set $H \subseteq A$ or $H \subseteq \bar{A}$ whose jump is not of DNC degree relative to \emptyset' ?

CONCLUSION

The computability-theoretic framework about Ramsey's theorem seems to be the good one.

There is hope to answer remaining questions about Ramsey's theorem using the standard combinatorics.

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